

Hands-on Breakout: Darshan

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Understanding I/O problems in your application

Example questions:

- □ How much of your run time is spent reading and writing files?
- Does it get better, worse, or is it the same as you scale up?
- Does it get better, worse, or is it the same across platforms?
- How should you prioritize I/O tuning to get the most bang for your buck?
- We recommend using a tool called Darshan as a starting point.
- In this hands-on session, we'll cover:
 - 1. Darshan background
 - 2. Darshan usage on HPC systems (e.g., ALCF Polaris)
 - 3. Darshan analysis tool insights
 - 4. General HPC I/O best practices and tuning considerations









What is Darshan?

Darshan is a scalable HPC I/O characterization tool. It captures a concise picture of application I/O behavior with minimal overhead.

★ Widely available

- Deployed at most large supercomputing sites
- Including most systems at ALCF, OLCF, and NERSC

★ Easy to use

- No changes to code or development process
- Negligible performance impact: just "leave it on"
- ★ Produces a *summary* of I/O activity for every job
 - This is a great starting point for understanding your application's data usage
 - Includes counters, timers, histograms, etc.



How does Darshan work?

Two primary components:

1. Darshan runtime library

- Instrumentation modules: lightweight wrappers (interposed at link or run time) intercept application I/O calls and record statistics about file accesses
 - File records are stored in bounded, compact memory on each process
- <u>Core library</u>: aggregate statistics when the application exits and generate a log file
 - Collect, filter, compress records and write a single summary file for the job



Figure courtesy Jakob Luettgau (UTK)



How does Darshan work?

Two primary components:

1. Darshan runtime library

<u>NOTE</u>: Though traditionally restricted to MPI apps, recent Darshan versions can often be made to work in non-MPI contexts. **More on this later...**



Figure courtesy Jakob Luettgau (UTK)



How does Darshan work?

Two primary components:

2. Darshan log analysis tools

- Tools and interfaces to inspect and interpret log data
 - PyDarshan command line utilities like the new job summary tool
 - Python APIs for usage in custom tools, Jupyter notebooks, etc.
 - Legacy C-based tools/library



Figure courtesy Jakob Luettgau (UTK)



Using Darshan

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Using Darshan

- We'll consider ALCF Polaris as an example in the following slides.
- The workshop repo includes Darshan hands on examples that are configured for use on Polaris.
 - <u>https://github.com/argonne-lcf/ALCF_Hands_on_HPC_Workshop</u>
 - See the darshan-hands-on directory
- Darshan deployments on other HPC systems are very similar. The most likely differences are:
 - Location of log files (where to find data after your job completes)
 - Analysis utility availability (usually easiest to just copy logs to your workstation to analyze)
 - Loading the Darshan module (if it's not already there by default)
- We'll briefly cover differences on other notable DOE systems after the Polaris example.



Using Darshan on Polaris: load the software

snyder@polaris-login-01> snyder@polaris-login-01> source polaris-setup-env.sh snyder@polaris-login-01>

snyder@polaris-login-01> cd ALCF Hands on HPC Workshop/darshan-hands-on/

The darshan-hands-on directory in the workshop GitHub repo includes a script to configure your environment with the tools needed for Darshan analysis.

<u>NOTE</u>: This additional setup script is manually loading the Darshan module, which is not yet enabled by default on Polaris – we are working on making this automatic!



Using Darshan on Polaris: load the software



The darshan-hands-on directory in the workshop GitHub repo includes a script to configure your environment with the tools needed for Darshan analysis.

Use "module list" to see a list of software loaded in your environment.

Darshan 3.4.3 should now be loaded.



Using Darshan on Polaris: load the software

snyder@polaris-login-01> cd ALCF_Hands_on_HPC_Workshop/darshan-hands-on/ snyder@polaris-login-01> snyder@polaris-login-01> source polaris-setup-env.sh snyder@polaris-login-01> snyder@polaris-login-01> module list

Currently Loaded Modules:

1)	craype-x86-rome	9)	cray-pmi/6.1.2
2)	libfabric/1.11.0.4.125	10)	cray-pmi-lib/6.0.17
3)	craype-network-ofi	11)	cray-pals/1.1.7
4)	perftools-base/22.05.0	12)	cray-libpals/1.1.7
5)	nvhpc/21.9	13)	PrgEnv-nvhpc/8.3.3
6)	craype/2.7.15	14)	craype-accel-nvidia80
7)	cray-dsmml/0.2.2	15)	darshan/3.4.3
8)	crav-mpich/8.1.16	16)	cray-python/3.9.12.1

These steps are similar, and often cases easier, on other HPC platforms where Darshan is deployed:

- Theta/Summit: Darshan module loaded by default
- Perlmutter: Darshan can be manually loaded with 'module load darshan'

Always check facility documentation!



Using Darshan on Polaris: instrument your code

Compile and run your application!

cc -o helloworld helloworld.c

qsub helloworld.qsub

From the helloworld example in the darshan-hands-on directory in the workshop GitHub repo

That's all there is to it; Darshan does the rest.*

* Well, almost. There is one caveat: in the default Darshan configuration, your application must call MPI_Init() and MPI_Finalize() to generate a log.



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Using Darshan on Polaris: find your log file

All Darshan logs are placed in a central location. The 'darshan-config --log-path' command will provide the log directory location.

snyder@polaris-login-04:~> darshan-config --log-path /lus/grand/logs/darshan/polaris snyder@polaris-login-04:~> snyder@polaris-login-04:~> cd /lus/grand/logs/darshan/polaris/2023/10/10/

Go to subdirectory for the year / month / day your job executed.

Be aware of time zone (or just check adjacent days)! Polaris, for example, uses the GMT time zone and will roll over to the next day at 7pm local time.



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snyder@polaris-login-04:~> darshan-config --log-path
/lus/grand/logs/darshan/polaris
snyder@polaris-login-04:~>
snyder@polaris-login-04:~> cd /lus/grand/logs/darshan/polaris/2023/10/10/
snyder@polaris-login-04:/lus/grand/logs/darshan/polaris/2023/10/10> ls | grep snyder
snyder@polaris-login-04:/lus/grand/logs/darshan/polaris/2023/10/10> ls | grep snyder
snyder_helloworld_id1126258-60233_10-10-72928-17597401521854260646_1.darshan

File name includes your username, app name, and job ID.

For convenience, users often copy logs somewhere else to save/analyze.



Using Darshan on Polaris: analyze log

After locating your log, users can utilize Darshan log analysis tools for gaining insights into application I/O behavior. PyDarshan tools likely aren't available everywhere, but traditional tools like darshan-parser should be.

shane@s	hane-	x1-carbon:	darshan-parser	./log.darsha	n grep	POSI	X_BYTES_WRIT	TEN	sort -nr -k 5
POSIX	387	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413869452		/projects/radix
POSIX	452	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413865644		/projects/radix
POSIX	197	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413857652		/projects/radix
POSIX	5	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413852168		/projects/radix
POSIX	451	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413844532		/projects/radix
POSIX	64	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413823236		/projects/radix
POSIX	68	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413788992		/projects/radix
POSIX	195	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413663132		/projects/radix
POSIX	323	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413658668		/projects/radix
POSIX	132	6966057	7185861764086	POSIX_BYTE	S_WRITTE	N	5413648628		/projects/radix

If you know what you're looking for, darshan-parser can be a quick way to extract important I/O details from a log, e.g., the 10 most heavily written files, but it is not super user friendly...



Using Darshan on Polaris: generate summary report

The Polaris environment setup script in the darshan-hands-on directory in the workshop GitHub repo also enables support for PyDarshan analysis tools.

Generate an HTML summary report with PyDarshan using the following command: '**python -m darshan summary <log_path>**'.

snyder@polaris-login-04> source ~/ALCF_Hands_on_HPC_Workshop/darshan-hands-on polaris-setup-env.sh snyder@polaris-login-04> python -m darshan summary helloworld.darshan /opt/cray/pe/python/3.9.12.1/ttb/python5.9/stte-packages/scipy/__tntt__.py:138: UserWarning: A NumPy version >=1.16.5 a nd <1.23.0 is required for this version of SciPy (detected version 1.25.1) warnings.warn(f"A NumPy version >={np_minversion} and <{np_maxversion} is required for this version of " Report generated successfully. Saving report at location: /home/snyder/ALCF_Hands_on_HPC_Workshop/darshan-hands-on/helloworld/helloworld_report.html



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Generate an HTML summary report with PyDarshan using the following command: 'python -m darshan summary <log_path>'.

snyder@polaris-login-04> source ~/ALCF_Hands_on_HPC_Workshop/darshan-hands-on/polaris-setup-env.sh
snyder@polaris-login-04>
snyder@polaris-login-04> python -m darshan summary helloworld.darshan
/opt/cray/pe/python/3.9.12.1/lib/python3.9/site-packages/scipy/__init__.py:138: UserWarning: A NumPy version >=1.16.5 a
nd <1.23.0 is required for this version of SciPy (detected version 1.25.1)
warnings.warn(f"A NumPy version >={np minversion} and <{np maxversion} is required for this version of "
Report generated successfully.
Saving report at location: /home/snyder/ALCF_Hands_on_HPC_Workshop/darshan-hands-on/helloworld/helloworld_report.html</pre>

If successful, the tool should generate an HTML report matching the input log file name.

To analyze, it's likely easiest to copy the report to your own workstation to view in a browser.

Using Darshan on Polaris: generate summary report

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Generate an HTML summary report with PyDarshan using the following command: 'python -m darshan summary <log_path>'.

snyder@polaris-login-04> source ~/ALCF_Hands_on_HPC_Workshop/darshan-hands-on/polaris-setup-env.sh
snyder@polaris-login-04> python -m_darshan_summary_helloworld_darshan
/opt/cray/pe/python/3.9.12.1/lib/python3.9/site-packages/scipy/__init__.py:138: UserWarning: A NumPy version >=1.16.5 a
nd <1.23.0 is required for this version of SciPy (detected version 1.25.1)
warnings.warn(f"A NumPy version >={np_minversion} and <{np_maxversion} is required for this version of "
Report generated successfully.
Saving report at location: /home/snyder/ALCF Hands on HPC Workshop/darshan-hands-on/helloworld/helloworld report.html</pre>

<u>NOTE</u>: Ignore these Python warnings about version requirements, they should not cause any issues with report generation



What about other HPC systems?

• **Perlmutter** (NERSC):

- How to enable: 'module load darshan'
- Log directory: /pscratch/darshanlogs/
- Summit (OLCF):
 - How to enable: automatic
 - Log directory: /gpfs/alpine/darshan/summit

If Darshan is not available on a system, it can be installed via Spack or directly from source. Darshan is provided as 2 separate packages in Spack:

- **darshan-runtime** library for instrumenting apps
- **darshan-util** tools for analyzing Darshan log files

PyDarshan is available on PyPI (e.g., '**pip install darshan**') and also in Spack

See our website for more details: https://www.mcs.anl.gov/research/projects/darshan



Analyzing Darshan logs



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Job analysis example

The PyDarshan job summary tool generates an HTML report containing graphs, tables, and performance estimates characterizing the I/O workload of the application

We will summarize some of the highlights in the following slides

e3sm_io (2022-03-02) **Darshan Summary Report** Job Summary Job ID 12345 6789 User ID 512 # Processes Run time (s 727.0000 2022-03-02 13:52:46 Start Time 2022-03-02 14:04:52 End Tin ./E3SM-IO/build/src/e3sm_io E3SM-IO-inputs/i_case_1344p.nc -k -o can_I_out.nc -a pnetcdf -x canonical -r 200 **Darshan Log Information** Log Filename e3sm_io_heatmap_only.darshan **Runtime Library Version** 3.3.1 Log Format Version 3.21 POSIX (ver=4) Module Data 20.34 KiB MPI-IO (ver=3) Module Data 0.49 KiB PNETCDE FILE (ver=2) Module Data 0 11 KiB LUSTRE (ver=1) Module Data 9.36 KiB STDIO (ver=2) Module Data 0.08 KiB APMPI (ver=1) Module Data 136 94 KiB HEATMAP (ver=1) Module Data 318.85 KiB I/O Summary Heat Map: HEATMAP MPIIO Heat Map: HEATMAP POSIX 51 306 ≠ 204 102 r- Time bins: 114 184 Time (s) 545



Job analysis: high-level job info





Job analysis: I/O heatmaps



Heatmaps showcase application I/O intensity (r+w volume) across time, ranks, and interfaces – helpful for identifying hot spots, I/O and compute phases, etc.



Job analysis: I/O heatmaps



Heatmaps showcase application I/O intensity (r+w volume) across time, ranks, and interfaces – helpful for identifying hot spots, I/O and compute phases, etc.

Sum rank over time slices



Job analysis: I/O heatmaps



This application demonstrates a couple of notable I/O characteristics:

- I/O imbalance across MPI processes
- Collective MPI-IO accesses transformed to subset of "aggregator" ranks at POSIX level



Job analysis: I/O cost



I/O cost indicates how much time on average was spent reading, writing, and doing metadata across different I/O interfaces

If I/O cost is a small portion of application runtime, tuning efforts are likely to have a relatively small impact



Job analysis: Per-interface statistics





Job analysis: Per-interface statistics



Operation counts provide the relative totals of different types of I/O operations

Lots of metadata operations (open, stat, seek, etc.) could be a sign of poorly performing I/O

Access pattern indicates whether read/write operations progress sequentially or consecutively* through the file

More random access patterns can be expensive for some types of storage



Job analysis: Per-interface statistics



Details on access sizes are provided to better understand granularity of application read/write accesses

In general, larger access sizes (e.g., O(MiBs)) perform better with most storage systems



Job analysis: Data access by category



Data accesses, in terms of total files read/written and total bytes read/written, binned by different categories:

- FS mount points (e.g., /home, /scratch)
- standard streams (e.g., STDOUT)
- object storage pools
- etc.

Inform on job's general usage of different storage resources



Job analysis: additional help

e3sm_io (2022-03-02)

Darshan Summary Report

Job Summ	hary
Job ID	12345
User ID	6789
# Processes	512
Run time (s)	727.0000
Start Time	2022-03-02 13:52:46
End Time	2022-03-02 14:04:52
Command Line	./E3SM-I0/build/src/e3sm_io E3SM-I0-inputs/i_case_1344p.nc -k -o can_l_out.nc -a pnetcdf -x canonical -r 200

Darshan Log Information

Log Filename	e3sm_io_heatmap_only.darshan
Runtime Library Version	3.3.1
Log Format Version	3.21
POSIX (ver=4) Module Data	20.34 KiB
MPI-IO (ver=3) Module Data	0.49 KiB
PNETCDF_FILE (ver=2) Module Data	0.11 KIB
LUSTRE (ver=1) Module Data	9.36 KIB
STDIO (ver=2) Module Data	0.08 KIB
APMPI (ver=1) Module Data	136.94 KiB
HEATMAP (ver=1) Module Data	318.85 KiB

I/O Summary







Remember to contact facility support staff for help! The Darshan job summary can be a good discussion starter if you aren't sure how to proceed with performance tuning or problem solving.



Darshan: a quick recap

- These slides have thus far covered some basic Darshan usage and tips.
- Key takeaways:
 - Tools are available to help you understand how your application accesses data.
 - The simplest starting point is Darshan.
 - It's likely already instrumenting your application, or can quickly be made to do so.
 - You will probably start with an HTML report generated using PyDarshan.
- Refer to documentation and support channels provided by the Darshan team and/or facilities staff.
 - <u>darshan-io.slack.com</u>





Hands-on Intermission 1

Are there any initial questions/comments about Darshan or how to use it?

As a starting point, you can try running an example program with Darshan instrumentation enabled to ensure the toolchain works as expected

- Follow instructions at darshan-hands-on/README.md to setup environment, compile and run examples, find Darshan output, and run analysis tools
 - You'll have to copy generated HTML reports to your own workstation to view in a browser using scp command
- Use the **helloworld** example or try it with an application of your own
- □ How many files did the application open?
- □ How much data did it read, and how much data did it write?
- □ What approximate I/O performance did it achieve?



HPC I/O insights with Darshan



A quick survey of the HPC I/O landscape

A complex data management ecosystem

As we discussed this morning, the HPC I/O landscape is deep and vast

- High-level data abstractions: HDF5, PnetCDF
- <u>I/O middleware</u>: MPI-IO
- <u>Storage systems</u>: Lustre, GPFS, DAOS
- Storage hardware: HDDs, SSDs, SCM

HPC applications themselves are evolving and encountering new data management challenges

Understanding I/O behavior in this environment is difficult, much less turning observations into actionable I/O tuning decisions





A quick survey of the HPC I/O landscape

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- <u>Storage systems</u>: Lustre, GPFS, DAOS
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Understanding I/O behavior in this environment is difficult, much less turning observations into actionable I/O tuning decisions

Darshan can help navigate this complexity by characterizing I/O behavior across the I/O stack




A look under the hood of an HPC application

You have already heard some basics about Darshan, a powerful tool for users to better understand and tune their I/O workloads

Darshan provides many helpful stats across multiple layers of the I/O stack that are critical to understanding application I/O behavior



HDF5 stats*:

- Accessed files/datasets
- Operation counts
- Total read/write volumes
- Common access info (including details of hyperslab accesses)
- Chunking parameters
- Dataset dimensionality and size
- MPI-IO usage
- I/O timing

*Note: HDF5 instrumentation is not typically enabled for facility Darshan installs – you will need to install this version yourself



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MPI-IO stats:

- Operation counts (open, read, write, sync)
- Collective and independent I/O usage
- Total read/write volumes
- Access size info
 - Common values
 - Histograms
- I/O timing

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POSIX stats:

- Operation counts (open, read, write, seek, stat)
- Total read/write volumes
- File alignment
- Access size/stride info
 - Common values
 - Histograms
- I/O timing

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Lustre stats:

- Data server (OST) and metadata server (MDT) counts
- Stripe size/width
- OST list serving a file



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I/O Hardware

Let's see how Darshan can be leveraged in some practical use cases that demonstrate some general best practices in tuning HPC I/O performance



Ensuring storage resources match application I/O needs

For some parallel file systems like Lustre, users have direct control over file striping parameters

Bad news: Users may have to have some knowledge of the file system to get good I/O performance

Good news: Users can often get higher I/O performance than system defaults with thoughtful tuning -- file systems aren't perfect for every workload!



Ensuring storage resources match application I/O needs

Tuning decisions can and should be made independently for different file types



Simulation clients write data to 1 storage server



Ensuring storage resources match application I/O needs

Tuning decisions can and should be made independently for different file types

Large application datasets should ideally be distributed across as many storage resources as possible



Simulation clients write data to 1 storage server



Simulation clients load balance writes across multiple servers



Ensuring storage resources match application I/O needs

Tuning decisions can and should be made independently for different file types On the other hand, smaller files often benefit from being stored on a single server



Simulation clients read config data from 1 storage server



Ensuring storage resources match application I/O needs

Tuning decisions can and should be made independently for different file types On the other hand, smaller files often benefit from being stored on a single server



using communication (e.g., MPI)



Ensuring storage resources match application I/O needs

Be aware of what file system settings are available to you and don't assume system defaults are always the best... you might be surprised what you find

 ALCF Polaris/Theta and NERSC Perlmutter Lustre scratch file systems both have a default stripe width of 1 (i.e., files are stored on one server by default)



256 process (4 node) h5bench¹ runs on NERSC Perlmutter

- 10⁹

107

106

10⁵ m

104 b

- 10³

10²

h5bench contains lots of parameters for controlling characteristics of generated HDF5 workloads.

1. https://github.com/hpc-io/h5bench



Ensuring storage resources match application I/O needs

Be aware of what file system settings are available to you and don't assume system defaults are always the best... you might be surprised what you find

 ALCF Polaris/Theta and NERSC Perlmutter Lustre scratch file systems both have a default stripe width of 1 (i.e., files are stored on one server by default)



All I/O funneled through rank 0

- 109

107

106

10⁵ m

104 pt

- 10³

10²

MPI-IO collective I/O driver for Lustre assigns dedicated aggregators for each stripe, yielding a single aggregator for files of 1 stripe



Ensuring storage resources match application I/O needs



MPI-IO



POSIX

Manually setting the stripe width to 16 yields more I/O aggregators and better performance:

> lfs setstripe -c 16 testFile



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Ensuring storage resources match application I/O needs



MPI-IO



POSIX

Manually setting the stripe width to 16 yields more I/O aggregators and better performance:

> lfs setstripe -c 16 testFile

4x performance improvement!



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Ensuring storage resources match application I/O needs

Consult facilities documentation for established best practice!

	Single Shared-File I/O	File per Process
File size	Command	Command
< 1 GB	keep default striping	keep default striping
1 - 10 GB	stripe_small	keep default striping
10 - 100 GB	stripe_medium	keep default striping
100 GB - 1 TB	stripe_large	keep default striping
> 1 TB	stripe_large	stripe_large

Perlmutter (NERSC) docs providing commands to set stripe params for various file types

- The default striping set on Orion is targeted to work well for a variety of workloads
- In most cases, users should use this default striping. Though possible, manual striping should only occur after careful consideration and under collaboration with OLCF staff
- The default striping policy may change due to findings in production

OLCF presentation on Orion storage system detailing usage of Lustre's new progressive file layout mechanism



Ensuring storage resources match application I/O needs

Consult facilities documentation for established best practice! Sometimes you may even need to experiment yourself.



128-node example of the IOR benchmark using various stripe counts on ALCF Polaris.

For more I/O intensive programs, it's typically better to err on the side of more storage servers. The following command stripes across all servers:

> lfs setstripe -c -1 testFile

https://github.com/radix-io/io-sleuthing/tree/main/examples/striping



Making efficient use of a no-frills I/O API

Users may also need to pay close attention to file system alignment when issuing I/O accesses to a file

• Accesses that are not aligned can introduce performance inefficiencies on file systems



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For Lustre, performance can be maximized by aligning I/O to stripe boundaries:



Unaligned I/O requests can span multiple servers and introduce inefficiencies in storage protocols



Making efficient use of a no-frills I/O API

Users may also need to pay close attention to file system alignment when issuing I/O accesses to a file

• Accesses that are not aligned can introduce performance inefficiencies on file systems

For Lustre, performance can be maximized by aligning I/O to stripe boundaries:



Instead, ensure client accesses are well-aligned to avoid Lustre server Fi contention





Making efficient use of a no-frills I/O API

Consider a simple 10-process (10-node) NERSC Cori example where processes write in an interleaved fashion to a single shared file

aligned

# Module	Rank	Wt/Rd	Segment	Offset	Length	Start(s)	End(s) [OST]
X_POSIX	0	write	Θ	0	1048576	0.0054	0.0066 [197]
X_POSIX	0	write	1	10485760	1048576	0.0066	0.0073 [197]
X_POSIX	0	write	2	20971520	1048576	0.0073	0.0081 [197]
X_POSIX	0	write	3	31457280	1048576	0.0081	0.0088 [197]

Use Darshan's DXT tracing module to get details about each individual write access – more details on DXT usage coming soon



Making efficient use of a no-frills I/O API

Consider a simple 10-process (10-node) NERSC Cori example where processes write in an interleaved fashion to a single shared file

aligned

# Module Rar	nk	Wt/Rd	Segment	Offset	Length	Start(s)	End(s)	[OST]
X_POSIX	0	write	0	Θ	1048576	0.0054	0.00	6 [197]
X_POSIX	0	write	1	10485760	1048576	0.0066	0.00	3 [197]
X_POSIX	0	write	2	20971520	1048576	0.0073	0.00	1 [197]
X_POSIX	0	write	3	31457280	1048576	0.0081	0.00	8 [197]

Each access is aligned to the Lustre stripe size (1 MiB)

Each process interacts with a single Lustre server (OST)



Making efficient use of a no-frills I/O API

Consider a simple 10-process (10-node) NERSC Cori example where processes write in an interleaved fashion to a single shared file

unaligned

# Module	Rank	Wt/Rd	Segment	Offset	Length	Start(s)	End(s) [OST]	
X_POSIX	0	write	0	524288	1048576	0.0065	0.05 4 [32	[197]
X_POSIX	0	write	1	11010048	1048576	0.0594	0.1208 [32	[197]
X_POSIX	0	write	2	21495808	1048576	0.1268	0.2000 [32	[197]
X_POSIX	0	write	3	31981568	1048576	0.2060	0.20(9 [32	[197]

Each access spans two Lustre stripes due to unaligned offsets

Each process interacts with two Lustre servers (OSTs)



Making efficient use of a no-frills I/O API

Even in this small workload, we pay a nearly **20% performance penalty when I/O accesses are not aligned** to file stripes (1 MB)

aligned

# Madula	Daak	LI+ /Dd	Compost	Offeet	Longth	C + 2 + (c)	[nd(c)	
# Module	RAIIK	WL/RU	Segment	Uliset	Length	Start(S)	End(S)	
X_POSIX	0	write	Θ	0	1048576	0.0054	0.0	310.14
X_POSIX	0	write	1	10485760	1048576	0.0066	0.0	
X_POSIX	0	write	2	20971520	1048576	0.0073	0.0	MiB/s
X POSIX	0	write	3	31457280	1048576	0.0081	0.0	

unaligned

# Module	Rank	Wt/Rd	Seament	Offset	Length	Start(s)	End(s) [05]	
X POSTX	0	write	0 0	524288	1048576	0.0065	0.0594	380.28
X POSTX	õ	write	1	11010048	1048576	0.0594	0.1268	300.20
X POSTX	õ	write	2	21495808	1048576	0.1268	0.2060	MiR/s
X POSIX	0	write	3	31981568	1048576	0.2060	0.2069	



Making efficient use of a no-frills I/O API

Accounting for subtle I/O performance factors like file alignment can be a painstaking process...

High-level I/O libraries like HDF5 can help mask much of the complexity needed for transforming scientific computing I/O workloads into performant POSIX-level file system accesses – **don't reinvent the wheel, use high-level I/O libraries wherever you can!**



Optimizing application interactions with the I/O stack

The HDF5 library provides a chunking mechanism to partition user datasets into contiguous chunks in the underlying file

 Users can greatly improve performance of partial dataset I/O operations by choosing chunking parameters that match expected access patterns



Optimizing application interactions with the I/O stack

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By default, HDF5 will store the dataset contiguously row-by-row (i.e., row-major format) in the file



Optimizing application interactions with the I/O stack

The HDF5 library provides a chunking mechanism to partition user datasets into contiguous chunks in the underlying file

 Users can greatly improve performance of partial dataset I/O operations by choosing chunking parameters that match expected access patterns



column-based



block-based

If dataset access patterns do not suit a simple row-major storage scheme, chunking can be applied to map chunks of dataset data to contiguous regions in the file



Optimizing application interactions with the I/O stack

Consider a 256-process (4-node) Polaris example where each process exclusively writes a 2048x2048 block of the dataset (32 MB per-process, 8 GB total)



With no chunking, each process must issue many smaller non-contiguous I/O requests (solid lines) and seek around the file (dashed lines), yielding low I/O performance



Optimizing application interactions with the I/O stack

Consider a 256-process (4-node) Polaris example where each process exclusively writes a 2048x2048 block of the dataset (32 MB per-process, 8 GB total)



256 individual HDF5 writes (1-per-process) yields 500K+ POSIX writes



Optimizing application interactions with the I/O stack

Consider a 256-process (4-node) Polaris example where each process exclusively writes a 2048x2048 block of the dataset (32 MB per-process, 8 GB total)



With chunking applied, each process can read their entire data block using one large, contiguous access in the file



Optimizing application interactions with the I/O stack

Consider a 256-process (4-node) Polaris example where each process exclusively writes a 2048x2048 block of the dataset (32 MB per-process, 8 GB total)



Chunking results in a much more manageable POSIX workload

> Nearly a 3x performance improvement!



Optimizing application interactions with the I/O stack

An alternative optimization forgoes chunking and uses collective I/O to improve the efficiency of this block-style data access

 Rely on MPI-IO layer collective buffering algorithm to generate contiguous storage accesses and to limit number of clients interacting with storage system



With collective I/O enabled, designated aggregator processes perform I/O on behalf of their peers, and communicate their data using MPI calls

E.g., the green process sends its write data to the blue process (aggregator), who then writes both of their data in one big contiguous chunk



Optimizing application interactions with the I/O stack

Consider a 256-process (4-node) Polaris example where each process exclusively writes a 2048x2048 block of the dataset (32 MB per-process, 8 GB total)





Collective I/O yields 26x improvement over no chunking, and 9x improvement over chunking!!!



Optimizing application interactions with the I/O stack

MPI-IO

252

Rank

252

198

153

99 -

54



POSIX

Darshan I/O activity heatmaps illustrate how different the I/O behavior is for the unoptimized independent configuration (top) and the most performant collective I/O configuration (**bottom**)



Summarizing I/O tuning options

As a user of I/O interface X, what tuning vectors do I have?

I/O Interface	Striping	Alignment	Collective I/O	Chunking
HDF5	✓	✓	✓	✓
MPI-IO	✓	✓	 Image: A second s	X
POSIX	✓	√ -	X	X



Summarizing I/O tuning options

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I/O Interface	Striping	Alignment	Collective I/C	D Chunking
HDF5	✓		 Image: A start of the start of	✓
MPI-IO	✓		1	X
POSIX	<pre> //</pre>		X	X
		Ţ		
Automatica data and libra re	lly align application ary metadata, if use quests so	collective l/ er be automated aligned	O can tically r d	POSIX I/O requires nanually aligning every access


Summarizing I/O tuning options

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I/O Interface	Striping	Alignment	Collective I/O	Chunking
HDF5	1	1	 Image: A start of the start of	1
MPI-IO		1	 Image: A second s	X
POSIX		√ -	X	X

Just another reminder that high-level I/O libraries are here to make your life easier

 I/O optimization strategies like collective I/O & chunking can net large performance gains, especially when combined with striping and alignment optimizations



Hands-on Intermission 2

In the darshan-hands-on directory there are 2 more examples with A & B versions: warpdrive & fidgetspinner

- Follow instructions at darshan-hands-on/README.md to setup environment, compile and run examples, find Darshan output, and run analysis tools
- Note: these examples will require at least some understanding of the MPI-IO library

See if you can spot the performance differences! Which version is faster? Why?

- Use Darshan job summary tool to compare I/O behavior
- Compare source code using diff to confirm



Additional Darshan tips and tricks



Darshan instrumentation beyond MPI

- Historically, Darshan has only worked with MPI applications
 - MPI_Init/MPI_Finalize used to bootstrap/shutdown Darshan
- Darshan has been modified to use a secondary bootstrapping mechanism that enables its use outside of MPI
 - Based on GCC-specific library constructor/destructor attributes
 - Only works for dynamically-linked executables!
- To enable non-MPI mode, users must explicitly opt-in by setting the DARSHAN_ENABLE_NONMPI environment variable
 - A unique log will be generated for every process that executes
 - Often best to limit instrumentation scope to the target executable:

```
$ LD_PRELOAD=/path/to/libdarshan.so \
    DARSHAN_ENABLE_NONMPI=1 \
    ./exe <args>
```





Finer-grained details with Darshan: DXT tracing

- By default, Darshan captures a fixed set of counters for each file
- With DXT, Darshan additionally traces every read/write operation (for POSIX and MPI-IO interfaces)
- Enable by setting **DXT_ENABLE_IO_TRACE** env variable
- Finer grained instrumentation data comes at a cost of additional overhead and larger logs

```
export DXT_ENABLE_IO_TRACE=1
mpiexec -n 256 --ppn 64 ./helloworld
```



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# #	DXT, DXT,	<pre>file_id: : rank: 0, 1</pre>	11637741 nostname	10118722858, : x3202c0s1b0	file_name: /g)n0	rand/projects/A	TPESC2023/us	r/snyder/hello
#	DXT,	write_cou	nt: 160,	read_count:	0			
#	DXT,	<pre>mnt_pt: /</pre>	, fs_typ	e: overlay				
#	Modu	le Rank	Wt/Rd	Segment	Offset	Length	Start(s)	End(s)
Х	POS	IX 0	write	Θ	Θ	1048576	3.9347	3.9468
Х	_POSI	IX 0	write	1	167772160	1048576	4.2503	4.2575
Х	POS	IX 0	write	2	335544320	1048576	4.5495	4.5564
Х	POS	IX 0	write	3	503316480	1048576	4.8632	4.8707

Trace includes the timestamp, file offset, and size of every I/O operation on every rank.

darshan-dxt-parser utility can provide a raw text dump of the trace



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Traces can be visualized using summary report heatmaps or other custom tools like DXT Explorer



Finer-grained details with Darshan: disabling shared file reductions

- To reduce log file size, globally shared file records are reduced into a single instrumentation record by default
 - However, this slightly masks per-rank contributions to I/O
- This behavior can be disabled by setting DARSHAN_DISABLE_SHARED_REDUCTION environment variable
- Allows for full accounting of per-rank contributions to shared files, if these details are important (e.g., for understanding collective I/O algorithms)

export DARSHAN_DISABLE_SHARED_REDUCTION=1

mpiexec -n 256 --ppn 64 ./helloworld \$SCRATCHDIR



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\$ darshan-parser ./snyder_helloworld_id565659-63984_8-3-68717-7310522192037150959_1.dar
> grep_POSIX_BYTES_WRITTEN
POSIX -1 1163774110118722858 POSIX_BYTES_WRITTEN 26214400000 /grand/

Rank -1 indicates a shared file record, with counters containing a reduced value access all ranks (e.g., ~24.5 GiB total bytes written across all ranks)



Finer-grained details with Darshan: disabling shared file reductions

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With shared reductions disabled, each rank retains their own record giving full insight into per-rank contributions (rank 0 writes 157 MiB and rank 255 writes nothing)



Darshan runtime library configuration

- To bound memory overheads, Darshan imposes several internal memory limits (total memory usage, per-module record limits, etc.)
- For some workloads, default limits may be exceeded resulting in partial instrumentation data
- To offer user's more control over memory limits and instrumentation scope, Darshan provides a comprehensive runtime configuration system
 - Environment variables or config files

# KEV	VALUE	MODULES
NAME_EXCLUDE	^/home	*
NAME_EXCLUDE	.pyc\$	*
NAME_EXCLUDE	.so\$	*
NAME_INCLUDE	.h5\$	*
MODMEM 8		
MAX_RECORDS	4000	POSIX
MOD_ENABLE	DXT_POSIX,DXT_MPIIO	
APP_EXCLUDE	git,ls,sed	

Regular expressions can be specified to control whether matching record name patterns are included/excluded in Darshan instrumentation



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# KEY NAME_EXCLUDE NAME_EXCLUDE NAME_EXCLUDE NAME_INCLUDE	VALUE ^/home .pyc\$.so\$.h5\$	MODULES * * * *
MODMEM 8		
MAX_RECORDS	4000	POSIX
MOD_ENABLE	DXT_POSIX,DXT_MPIIO	
APP_EXCLUDE	git,ls,sed	

Settings are also offered to control total per-process memory usage (8 MiB) and per-module maximum record counts (4000 POSIX records)



Darshan runtime library configuration

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# KEY NAME_EXCLUDE NAME_EXCLUDE NAME_EXCLUDE NAME_INCLUDE MODMEM 8	VALUE ^/home .pyc\$.so\$.h5\$	MODULES * * *
MAX_RECORDS	4000	POSIX
MOD_ENABLE	DXT_POSIX,DXT_MPIIC	
APP_EXCLUDE	git,ls,sed	

Additional settings allow control over enabled/disabled modules, as well as application names that should be included/excluded from instrumentation



Thank you!

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Bonus

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Darshan-based analysis tools

Using Darshan as a starting point for developing new I/O analysis tools is attractive for a couple of reasons:

- 1. Darshan is commonly deployed in production at many HPC sites, making its I/O characterization data generally accessible to custom tools
- 2. Recent PyDarshan work has enabled much more agile development of Darshan-based I/O analysis tools in Python

We will start by considering a couple of Darshan-based I/O analysis tools: **DXT Explorer** and **Drishti**



DXT Explorer

- Darshan does not offer much in terms of DXT trace analysis tools beyond general I/O activity heatmaps
- DXT Explorer is an interactive web-based trace analysis tool for DXT data that was developed to provide:
 - Combined views of MPI-IO and POSIX activity
 - Zoom in/out capabilities to focus on subsets of ranks or specific time slices
 - Contextual information about I/O calls
 - Views based on operation type, size, and spatiality
- Interactive trace analysis with DXT Explorer can enable interesting new insights into app I/O behavior



★ DXT Explorer was developed by Jean Luca Bez (LBL). Slide content also provided courtesy of Jean Luca.

Bez, Jean Luca, et al. "I/O bottleneck detection and tuning: connecting the dots using interactive log analysis." *2021 IEEE/ACM Sixth International Parallel Data Systems Workshop (PDSW)*. IEEE, 2021.





Explore the timeline by zooming in and out and observing how the MPI-IO calls are translated to the POSIX layer.

For instance, you can use this feature to detect stragglers.



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Explore the spatiality of accesses in file by each rank with contextual information.

Understand how each rank is accessing each file.



Drishti

- Darshan can capture detailed I/O characterization data for an app, but translating this raw data to actionable tuning feedback is a significant challenge
- Drishti is a command-line tool to guide end-users in optimizing I/O in their applications by detecting typical I/O performance pitfalls and providing a set of recommendations
- Drishti checks each given Darshan log against 30+ heuristic triggers for various I/O issues and suggests actions to take to resolve them
 - 4 levels of triggers: high, warning, ok, info



github.com/hpc-io/drishti-io



docker pull hpcio/drishti

★ Drishti was developed by Jean Luca Bez (LBL). Slide content also provided courtesy of Jean Luca.

Bez, Jean Luca, Hammad Ather, and Suren Byna. "Drishti: guiding end-users in the I/O optimization journey." 2022 IEEE/ACM International Parallel Data Systems Workshop (PDSW). IEEE, 2022.



Drishti

	Drishti
- DRISHTI v.0.3 -	
JOB: EXECUTABLE: DARSHAN: EXECUTION DATE: FILES: PROCESSES HINTS:	1190243 bin/8_benchmark_parallel jlbez_8_benchmark_parallel_id1190243_7-23-45631-11755726114084236527_1.darshan 2021-07-23 16:40:31+00:00 to 2021-07-23 16:40:32+00:00 (0.00 hours) 6 files (1 use STDIO, 2 use POSIX, 1 use MPI-IO) 64 romio_no_indep_rw=true cb_nodes=4
- 1 critical issu	es, 5 warnings, and 5 recommendations
- METADATA	
 Application i Application m Application m 	s read operation intensive (6.34% writes vs. 93.66% reads) ight have redundant read traffic (more data was read than the highest read offset) ight have redundant write traffic (more data was written than the highest write offset)
- OPERATIONS	
► Application i read/write reque ↔ 284 (36.98% ↔ Recommendat ↔ Consider ↔ Since the	<pre>ssues a high number (285) of small read requests (i.e., < 1MB) which represents 37.11% of all sts) small read requests are to "benchmark.h5" ions: buffering read operations into larger more contiguous ones appplication already uses MPI-IO, consider using collective I/O calls (e.g. MPI_File_read_all() or application already uses MPI-IO, consider open </pre>

Overall information about the Darshan log and execution

Number of critical issues, warning, and recommendations

Details on metadata and data operations

Critical issue and corresponding recommendation for benchmark.h5

